

An Efficient Scheme for Securing Data Warehouses in the Cloud by Reducing Overhead While Enforcing Data Privacy

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Abstract: In cloud computing schemes, Infrastructure as a Service represents a cloud-computing technology that delivers computing resources, networking, and storage to consumers on-demand, over the internet. It enables end customer or end users to upscale or downsize resources on an as-when needed basis, reducing the need for upscaling, up-front capital expenditures or unnecessary infrastructure. In this paper, we propose an efficient additive encryption scheme based on Shamir's secret sharing for securing data warehouses in the cloud that addresses the shortcomings of existing approaches by reducing overhead while still enforcing good data privacy.

1. Introduction

The continued advancement of information technology and data communications strengthens the exchange of highly sensitive medical information. electronic health systems are widely used, and many medical facilities rely on the transmission and receipt of medical information online and on local networks. Over the years, many security systems have been introduced sqto monitor patient privacy and ensure the safety of interchangeable medical data. Cryptography is one of the techniques that often provides security for eHealth systems [1-5].

Nowadays, data outsourcing scenarios tremendously grow with the advent of cloud computing that offers both cost savings and service benefits. One of the most noteworthy cloud outsourcing services is Database-as-a-Service, where individuals and organizations outsource data storage and management to a Cloud Service Provider (CSP) [6-17]. Naturally, such services allow outsourcing a DW and running OLAP queries. Yet, data outsourcing brings out privacy concerns since sensitive data are stored, maintained and processed by an external thirdparty that may not be fully trusted.

A typical solution to preserve data privacy is encrypting data locally be-fore sending them to an external server. Secure database management systems (SDBMSs) such as CryptDB implement cryptographic schemes [18-29]. Paillier's partially homomorphic encryption scheme is notably used in CryptDB to provide high security. However, it induces a

high storage and computation over-head. We propose a new Secure Secret Splitting Scheme (S4) that aims at replacing Paillier's scheme in systems such as CryptDB. S4 is based on the idea of secret sharing and is efficient both in terms of storage and computing, without sacrificing privacy too much.

CryptDB brings together powerful cryptographic tools to handle query processing on encrypted data without decryption. Encryption in CryptDB is like onion layers that store multiple ciphertexts, i.e., encrypted data, within each other. Each onion layer enables certain kind of query processing and a given security level provided by one encryption scheme [30-43]. For instance, order-preserving encryption (OPE) enables range queries and additive homomorphic encryption enables addition over encrypted data. Yet, CryptDB is not perfectly secured since schemes such as OPE reveal some statistical information about plaintext. MONOMI builds upon CryptDB to allow the execution of analytical work-loads over encrypted data outsourced to the cloud. MONOMI aims at improving CryptDB's query processing capability and efficiency based on split client/server execution. A designer also optimizes physical data layout [44-61]. Eventually, using a local trusted hardware at the CSP's, such as TrustedDB and CipherBase, is an alternative approach to query encrypted data. However, trusted hardware is limited in computation ability and memory capacity, and also very expensive.

Fully homomorphic encryption (FHE) allows performing arbitrary arithmetic operations over encrypted data without decryption. FHE provides semantic security, i.e., it is computationally impossible to distinguish two ciphertexts encrypted from the same plaintext. However, FHE requires so much computing power that it cannot be used in practice. Partially homomorphic encryption (PHE) is more efficient than FHE. Paillier's is the most efficient additive FHE. With Paillier's scheme, multiplying the encryption of two values results in an encryption of the sum of the values, i.e., $Enc_k(x) \cdot Enc_k(y) = Enc_k(x + y)$, where the multiplication is performed modulo some public-key k . Paillier's scheme is, however, still computationally intensive and induces as large ciphertext sizes as 2048 bits. Additionally, modular multiplications become computationally expensive on a large number of records, such as in the fact counter of a DW.

Secret sharing divides a secret piece of data into so-called shares that are stored at n participants'. A subset of k n participants is required to reconstruct the secret. In Shamir's, the first secret S4's driving idea is based on secret sharing, but instead of sharing secrets to n participants' or CSP's, they are stored at one single CSP's. Thus, we avoid the high storage

overhead of secret sharing. In S4, each secret v_j is divided into $n = k$ splits $v_{1,j}; \dots; v_{k,j}$. $k-1$ splits, $v_{1,j}; \dots; v_{k-1,j}$, are stored at the CSP's and $v_{k,j}$ is stored in a trusted machine, e.g., at the user's. In order to reduce storage overhead at the user's, $v_{k,j}$ is set to be the same for all secrets.

2. Methodology

First, x_k and v_k are randomly set up from F_p , where p is a big prime number, i.e., greater than the greatest possible query answer. For any secret v_j , a random polynomial $P_{v_j}(x)$ is built that passes through $(0; v_j)$ and $(x_k; v_k)$. To this end, $k-2$ points $(a_i; b_i); i = 1; \dots; k-2$ are chosen randomly from F_p such that $a_i \neq x_k$ and $a_i \neq 0 \forall i = 1; \dots; k-2$. Given k points $(a_1; b_1); (a_2; b_2); \dots; (a_{k-2}; b_{k-2}), (0; v_j)$ and $(x_k; v_k)$, polynomial $P_{v_j}(x)$ is built. Storing the $k-2$ random points is unnecessary because they are not needed for secret reconstruction. To divide v_j into $k-1$ splits (since $(x_k; v_k)$ is already fixed), a set of $k-1$ distinct elements $X = \{x_1; x_2; \dots; x_{k-1}\}$ is chosen from F_p such that $x_i \neq 0$ and $x_i \neq x_k \forall i = 1; \dots; k-1$. Then, splits are $v_{i,j} = P_{v_j}(x_i)$. $K = (X; (x_k; v_k))$ is considered as a private key for S4 and must be kept hidden from the CSP. To reconstruct secret v_j , its $k-1$ splits must be retrieved from the CSP. Given points $(x_i; v_{i,j}), i = 1; \dots; k-1$ and $(x_k; v_k)$, which is stored at the user's, polynomial $P_{v_j}(x)$ can be reconstructed.

Let a relational T consist of one attribute A (additional attributes, if any, can be processed similarly). Suppose T has m records. We denote by v_j the j^{th} value of A . For attribute A in T , $k-1$ attributes $A_i, i = 1; \dots; k-1$ are created in T^0 at the CSP's, where each attribute A_i stores the i^{th} splits. Without loss of generality, we assume integer data type for other data types can be transformed into integers before splitting. S4 allows summation queries to be computed directly at the CSP's. Consider a query that sums q values of A .

Paillier's PHE is semantically secure, but it is too expensive in terms of cipher-text storage space and query response time. S4 proposes a classical trade-off with a lower level of security, but better storage and response time efficiency. Let us consider a scenario where the CSP is said honest but curious, which is a widely used adversary model for cloud data outsourcing. Such a CSP faithfully complies to any service-level agreement and, in our particular case, stores data, runs queries and provides results without alteration, malicious or otherwise. Yet, the CSP may access data and infer information from queries and results.

Privacy in S4 relies on the fact that a secret value is only retrievable by the user via private key K . As in secret sharing, it is indeed guaranteed that at least k splits and X are necessary to reconstruct a secret, while the CSP has access to only $k-1$ splits. Both X and the k^{th} split, i.e., K , are stored at the user's. However, the CSP still has access to linear

combinations of splits, which provide some information. Still, the higher k is, the more difficult it is to interpret linear combinations of splits. Thus, k is the prime security parameter in S4. Experiments provide hints for choosing k .

Moreover, if some secrets are known by the CSP, e.g., through public communication of a company to its shareholders. For example, if the CSP knows secrets $v_1; \dots; v_{k-1}$. Also knowing the corresponding splits $v_{1,j}; \dots; v_{k-1,j}$ $\forall j \in \{1, \dots, k-1\}$, the CSP can recover the Lagrange basis polynomials $\ell_i(0) \forall i \in \{1, \dots, k-1\}$ and recover all secrets. However, the CSP must know at least $k-1$ secrets to do so. Moreover, we also propose leads to address this problem in next segment.

3. Results

We implement S4 in C using compiler gcc 4.8.2. S4's source code is freely available online. Experiments related to Paillier's PHE exploit the libpaillier standard C library. All mathematical computations use the GNU Multiple Precision Arithmetic Library (GMP). Eventually, we conduct our experiments on an Intel Core i7 3.10 GHz PC with 16 GB of RAM running Linux Ubuntu 15.05. We compare S4 and Paillier's PHE using simple synthetic datasets, i.e., 32-bit unsigned integers generated uniformly at random from the integer range $(103; 104)$. We scale up the number of records m such that $m \in \{103, 104, 105, 106\}$, forming four distinct datasets. In S4, we vary k from 8 to 64, higher values of k inducing too long execution times. Prime p must be greater than the greatest query answer, e.g., $p > \sum_{j=1}^m v_j$. In Paillier's PHE, we use a key size of 1024 bits, which induces ciphertexts of 2048 bits. Such key size is the absolute minimum to achieve security.

It is seen that encryption time in S4 is lower than Paillier's when $k \leq 16$, and then becomes higher when $k > 16$. Secret splitting consists in building a random polynomial by randomly choosing $k-2$ points. Hence, splitting time increases with k . This actually illustrates the tradeoff between S4's security and encryption efficiency with respect to Paillier's PHE. With the selected values of k , decryption is faster with S4 than with Paillier's PHE. This is mainly because Paillier's scheme needs m expensive modular multiplications of large, 2048-bit numbers for decryption, while secret reconstruction in S4 works by polynomial interpolation over k points and evaluating the polynomial in one single point.

With the selected values of k , S4's storage overhead is always much smaller than that of Paillier's PHE since y axis follows a logarithmic scale. Paillier's scheme indeed produces 2048-bit cipher texts. Thus, its storage overhead is $m \cdot 2048$. With S4, each value is split into k

1 values. Thus, S4's storage overhead is $m(k-1)$ times plaintext size. It is seen that, with the selected values of k , query execution time in S4 is lower than that of Paillier's scheme. This is because Paillier's scheme requires m expensive modular multiplications to compute a sum, while S4 computes only $(k-1)m$ simple modular additions.

4. Conclusions

We achieved performance gains through a slight degradation of security, especially when an adversary has knowledge of secret values. Although it is definitely satisfactory in some cloud DW and OLAP scenarios, e.g., public aggregate data might not actually yield secrets, i.e., fine-grained data, we will devote future research to strengthen S4 against such threats. More precisely, we plan to introduce noise, as in many cryptographic problems such as approximate-GCD or LWE.

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